

The problem considered in this monograph may be described as one of finding useful graphical descriptions of single contextual relations germane to a system. The problem is addressed as one of developing a process that minimizes the logistical burden on the developer of the product. The product is called an interpretive structural model. The latter is structurally isomorphic to a digraph, but contains both substantive and structural information supplied by the developer. All of the mathematical operations required by the process would be carried out by a computer. The critical, substantive system knowledge is supplied by the developer. The partnership of developer and computer gives the process its strength.

By a system is understood a complex of entities S with interactions R . Knowledge concerning (S, R) can be structural or substantive or both, and is the consequence of perceptions of the system. Some knowledge of the system is presumed to be available. Total knowledge of a system is thought normally to be unattainable. A description of a system is approximate.

Let $\langle S \rangle$ be represented by a finite set $S = \{s_1, s_2, \dots, s_n\}$ and let $\langle R \rangle$ be represented by a finite set $\langle R \rangle = \{\langle R_1 \rangle, \langle R_2 \rangle, \dots, \langle R_m \rangle\}$. Call s_i an element of the system, and call $\langle R_j \rangle$ a contextual relation of the system. A contextual relation is a phrase in a colloquial language.

Let $\langle R_j \rangle$, be represented by R_j , where R_j is a binary relation on $S \times S$. Let A_{ij} be a binary matrix indexed by (S, S) which portrays both R_j and its complement, not R_j . Let M_{ij} be the transitive closure of $A_{ij} + I$, where I is the identity matrix.

Then, there exists a permutation matrix P such that the matrix $M_j = (1/P)M_{ij}P$ induces a unique partition on S into subsets called levels, and each submatrix of M_j indexed by a level is symmetric.

For any submatrix indexed by a level, there exists a permutation such that its block diagonal submatrices form a minimum set of universal matrices. Each nontrivial index set of one or more of these is called a maximal cycle set. If an index set contains only one element, the set is trivial and the element is an isolate within the level, though it may connect to other levels.

If appropriate modifications are made to M_j , it can be converted to a matrix $M'_{j'}$ such that every maximal cycle set is replaced by a proxy element, and reachability is preserved by following a principle of condensation.

The digraph of the condensed matrix is a hierarchy with the same number of levels as M_{ij} and M_j . There exists a binary matrix A_j for which $M'_{j'}$ is the transitive closure, called a skeleton matrix, such that there is no other matrix less than A_j that has $M'_{j'}$ as its transitive closure. The digraph $D(A_j)$ is called the skeleton digraph of M_{ij} . It is the minimum-edge digraph that preserves condensed reachability or A_{ij} .

To every maximal cycle set, there corresponds a minimum-edge digraph that preserves reachability within the cycle set. If every proxy vertex on $D(A_j)$ is replaced with such a minimum edge digraph, the resulting digraph is a minimum-edge approximation of $D(A_{ij})$.

Because the matrix A_{ij} is often difficult to develop, the matrix M_{ij} is developed instead. In developing M_{ij} transitivity is enforced. While it is impossible to determine A_{ij} from M_{ij} , it is possible to determine A_j .

Since a minimum-edge representation of a maximal cycle is often an inadequate approximation to the structure of a maximal cycle, additional information concerning the structure of the cycle may be sought. Such structure is based upon a weighting matrix W for the maximal cycle, wherein the intensity of some contextual relation is supplied on a specified scale. From such a weighting matrix, it is possible to compute a binary matrix that represents a threshold of intensity and preserves reachability within the maximal cycle. This binary matrix permits construction of a reachability-preserving digraph for the maximal cycle which normally has fewer edges than a complete graph for the cycle.

If those digraphs that result from application of an intensity threshold to weighting matrices of the various maximal cycles are substituted for the proxy elements in the skeleton digraph $D(A_{ij})$ while preserving overall reachability, it is believed a good approximation to the system structure results insofar as the relation R_j is concerned. The digraph so formed may be complicated by the presence of many vertexes, edges, and crossings of the edges. If so, additional interpretive digraphs may be useful.

A maximal geodetic cycle specified on (s_i, s_j) is a maximal cycle consisting of a pair of (geodetic) paths, one path originating at vertex s_i and terminating at s_j , the other originating at s_j and terminating at s_i . It is possible to compute, for a given matrix A_{ij} , a set of maximal geodetic cycles. To each maximal geodetic cycle there corresponds a cycle set, which is a subset of S . The set of all such cycle